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Scalable Reflection in C++

Revision history

R0	Initial revision introducing scalar reflection model, reifiers (now called <i>splicers</i>), extensive API, and many examples.
R1	Added <code>refexpr(... xyz)</code> . Revised reifier/splicer syntax somewhat. Report on implementations. Introduce <code><meta></code> header. Reorganized presentation slightly.
R2	Harmonized with other papers in this area, including the use of the term “splicing” instead of “reifying” and the syntax developed in P2320R0. Various fixes and presentation improvements.

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Introduction

The first Reflection TS (based on N4766) exposes reflection information as types (to simplify integration with template metaprogramming techniques). However, SG7 agreed some time ago that the future of reflective constructs in C++ should be value-based (see also P0425r0). Specifically, the compile-time computations required for reflective metaprogramming should make use of constant-evaluation, which, unlike template metaprogramming, allows for ephemeral intermediate results (i.e., they don’t persist throughout the compilation process) and for mutable values. This approach was described in P0993r0, *Value-based Reflection*. To support that reflection design, we have passed a number of “`constexpr`” extensions in C++20: `consteval` functions (P1073), `std::is_constant_evaluated()` (P0595), and `constexpr` dynamic allocation (P0784), amongst others. We have also proposed *expansion statements* (P1306), which are more broadly useful but especially convenient for reflective metaprogramming: That feature was approved by the evolution working group for C++20, but did not make it to a WG21 vote for lack of time completing the Core wording review. We still hope expansion statements will be added to the language in the relatively near future.

That in itself still leaves plenty of design options for the reflection interface itself. What follows is an extensive document describing:

- The representation and properties of “reflections” (with argumentation for our specific design and considerations of alternatives).
- Mechanisms for *splicing*: Turning reflections into ordinary C++ source constructs (again, with design discussions).
- A brief discussion about templates and their instances.
- Principles to translate existing standard template metaprogramming facilities to the reflection domain.
- Principles to translate the Reflection TS facilities to the value-based reflection domain.
- Some examples to argue that proposals to add additional template metaprogramming facilities are unneeded because the underlying functionality is better handled in the reflection domain.
- An appendix listing the meta-functions being worked on one ongoing implementation.

This paper doesn’t exist in a vacuum. Related topics have been separately explored in P2320R0 (“The Syntax of Static Reflection”), P2237 (“Metaprogramming”), P2050R0 (“Tweaks to the design of source code fragments”), P1717 (“Compile-time Metaprogramming in C++”), and P1306 (“Expansion statements”).

Earlier versions of this paper were more exploratory in nature; this version uses experience with implementations based on earlier versions to narrow down a first set of metaprogramming features that are primarily aimed at providing reflection facilities (with splicing and ordinary template instantiation handling generative programming). However, additional facilities (particularly, for code injection) have been explored along with this proposal and we are not confident that they can be added incrementally on top of this proposal.

A simple example

The following function uses static reflection facilities presented in this paper to compute the string representation of an enumerator value.

```
#include <meta>
template<Enum T>
std::string to_string(T value) { // Could also be marked constexpr
    template for (constexpr auto e : std::meta::members_of(^T)) {
        if ([:e:] == value) {
            return std::string(std::meta::name_of(e));
        }
    }
    return "<unnamed>";
}
```

In broad strokes, the function does the following:

1. Gets the sequence enumerators from the enumeration type `T`,
2. Iterates over those enumerators, searching for the first that matches `value`,
3. Returns the name of that iterator.

Each of these operations relies on a feature included in this proposal. In particular, getting the sequence of iterators requires that we first get a queryable representation of the enumeration type `T`. This is done using the prefix `^` operator; it returns a *reflection*: a handle to an internal representation of type `T` maintained by the compiler. The `members_of` function (declared in a newly proposed standard header `<meta>`) returns a compile-time `std::span`, whose elements are the reflections of each enumerator in the enum.

To iterate over the span we use an *expansion-statement* (proposed through a separate paper P1306, and previously approved by EWG but still in CWG review), spelled `template for`. This isn't true "iteration", however. The body of the statement is repeated for each element in the span `s` so that the loop variable (`e` above) is initialized to `s[0], s[1], ..., s[n - 1]` in each successive repetition. The expansion variable is declared `constexpr` and carries into each repeated body. In other words, each repetition is equivalent to:

```
{
    constexpr std::meta::info e = s[I];
    if ([:e:] == value)
        return std::meta::name_of(e);
}
```

where `I` counts the repetitions of the loop's body.

Within the expansion body, the `[: ref1 :]` construct recovers the value of a reflected entity. We call this recovery process *splicing* and the constructs — like `[:...:]` — that enable it *splicers*. This can be compared with the parameter `value` to determine if they are the same. Finally, the `name_of` function returns a compile-time `string_view` for the identifier spelling of the matched enumerator. If none of the enumerators matched (possible, e.g., when bit-ORing together enumerator values), we return a string "`<unnamed>`" (which won't collide with a valid identifier).

This is called *static reflection* because all of the operations used to query types and enumerators are computed at compile time (i.e., statically). There is no additional runtime meta-information that must be generated with such facilities, which reinforces the zero-overhead principle that is so fundamental to C++. There is no runtime representation of the enumeration type and its enumerators. Only information that is ODR-used is present in the final program.

Implementation status

Two implementations of this proposal are underway.

The first and most complete is a fork of Clang by Lock3 Software (by, among others, Andrew and Wyatt, authors of this paper). It includes a large portion of the capabilities presented here, albeit not always with the exact syntax or interfaces proposed. In addition to these capabilities, Lock3's implementation supports expansion statements and injection primitives (including "fragment" support). Lock3 is currently not maintaining this implementation, however.

The second is based on the EDG front end (by Faisal and Daveed) and is less complete: It implements the reflection operator and most single splicers (but not the pack splicers; see below), and a few meta-library interfaces. It does not currently implement features in other proposals like expansion statements or injection primitives.

Reflections

The `^` operator

The first Reflection TS introduced the `reflexpr` operator to obtain reflection values encoded as types. Previous versions of this paper attempted to avoid repeating the considerable bikeshedding that went into selecting the `reflexpr` keyword by simply reusing it. Ironically, the spelling is more appropriate for the value-based reflection since the corresponding operation is indeed an "expression" (i.e., a construct that produces a value; in the TS it produces a type).

However, with months of practice with implementations that used `reflexpr(...)` we experienced consistent feedback that that syntax is too "heavy". So we went back to the drawing board and found that

the `^` prefix operator — suggesting “lifting” or “raising” representation — is available¹. This new syntax was agreed to by SG-7 during the discussion of P2320. Thus, we can write:

```
constexpr std::meta::info reflection = ^name_or_postfix_expr;
```

The value of `reflection` (i.e. the result of this *lifting operator*) is a compile-time value that *designates* some view of the indicated program element by the implementation (specifically, the compiler front end). I.e., it can be thought of as a handle to an internal structure of the compiler. In the rest of this proposal, we refer to the result of `^` as a *reflection* or a *reflection value*.

Note that the lifting operator is the “gateway” into the reflected world, but it is not the only source of *reflections* (or *reflection values*): We will further introduce a variety of functions that derive reflections from other reflections (e.g., we’ll present a function that returns reflections for the members of a class given a reflection for that class). Whatever the source of a reflection, we say that it *designates* language concepts such as entities or value categories. As will be shown later, a reflection can *designate* multiple notions. For example, `^f(x)` designates the called function `f` (if indeed that is what is called) and the type and value category of the call result.

The operand of `^` must be one of the following:

- a *type-id*, including possibly a *simple-type-specifier* that designates a *template-name*
- a possibly qualified *namespace-name*
- the scope-qualifier token “`:`” (designating the global namespace)
- a *postfix-expression*²

In the case where the `name_or_postfix_expr` is an expression, it is unevaluated but *potentially constant evaluated*. That implies that given “`struct S { int x; };`”, the expression “`^S::x`” is permissible in this context. We will elaborate the available reflected semantics later in this paper. Since `^name_or_postfix_expr` is an expression, `^(^name_or_postfix_expr)` is also valid (generally producing a distinct reflection).

In this paper, we call *declared entity* any of the following: a namespace (but not a namespace alias), a function or member function (that includes implicit special members, but not inherited constructors), a function or template parameter, a variable, a type (but not a type alias), a data member, a base class, a capture, or a template (including an alias template, but not a deduction guide template). Note that this is slightly different from the standard term *entity* (which, e.g., includes “values” but not “captures”³). We call *alias* a namespace alias or a type alias.

¹ We used to think that C++/CLI had already appropriated that syntax, but C++/CLI (and related C++ dialects) only uses the caret for handle declarations and *not* for handle indirections.

² Which includes any parenthesized expression.

³ This paper does not currently deal with structured bindings because their exact nature in the standard is still somewhat in flux at the time of this writing. Once they’re clarified, we intend to revisit their status as a “declared entity”.

Reflection type

What should the type of a reflection be? We propose it to be a new scalar type⁴, distinct from all other scalar types, that supports — aside from reading, assigning, and copying — only the scalar operations of equality/inequality and contextual conversion to `bool`. In addition we propose specific *splicers* (that transform a reflection value into a type or a name, see below) and library functions that can operate on `constexpr` reflections and `constexpr` sequences of reflections and generate new reflection-values as needed. All other operations on reflection values are then composed from these aforementioned operations. We present our rationale below for this design choice.

It is tempting to organize reflection values as class type values using a hierarchy of class types that try to model the language constructs. For example, one could imagine a base class `Reflection`, from which we might derive a class `ReflectedDeclaration`, itself the base class of `ReflectedFunction` and `ReflectedVariable`.

We do not believe that is the best approach for at least the following reasons:

- Although the relationship between major language concepts is relatively stable, we do occasionally make fundamental changes to our vocabulary (e.g., during the C++11 cycle we changed the definition of “variable”). Such a vocabulary change is more disruptive to a class hierarchy design than it is to certain other kinds of interfaces (we are thinking of function-based interfaces here).
- Class hierarchy values aren’t friendly to value-based programming because of slicing; instead, it works better with “reference” programming, which is particularly expensive for constant evaluation (because it requires address computations, which involve additional bookkeeping to check for potential undefined behavior).
- Class types are not easily used as nontype template arguments, particularly when we want to restrict effects to compile time (the recently added support for nontype class-type template arguments (P0732R2+P1907R1) imposes draconian limitations on class types). As it turns out, instantiating templates over reflection values is an important idiom when it comes to generative programming (e.g., through splicers or, eventually, code injection).
- Implementations of constant-evaluation usually handle non-pointer scalar values significantly more efficiently than class values.

Regarding this last point, the following compile-time test:

⁴ We could define it via `using info = decltype(^void)`.

```

constexpr int f() {
    int i = 0;
    for (int k = 0; k<10000; ++k) {
        i += k;
    }
    return i/10000;
}
template<int N> struct S {
    static constexpr int sm = S<N-1>::sm+f();
};
template<> struct S<0> {
    static constexpr int sm = 0;
};
constexpr int r = S<200>::sm;

```

compiles in about 0.6 seconds on a compact laptop (2016 MacBook m7), but wrapping the integers as follows:

```

struct Int { int v; };
constexpr int f() {
    Int i = {0};
    for (Int k = {0}; k.v<10000; ++k.v) {
        i.v += k.v;
    }
    return i.v/10000;
}
template<int N> struct S {
    static constexpr int sm = S<N-1>::sm+f();
};
template<> struct S<0> {
    static constexpr int sm = 0;
};
constexpr int r = S<200>::sm;

```

doubles the compile time to 1.2 seconds. Adding a derived-class layer would further increase the time. Another increase would result from attempting to access the classes through references (as would be tempting with a class hierarchy) because address computations require some work to guard against undefined behavior.

Because of these various considerations, we therefore propose that the type of a reflection is an unspecified scalar type, distinct from all other scalar types, whose definition is:

```
namespace std::meta {
    using info = decltype(^void);
}
```

Namespace `std::meta` is an associated namespace of `std::meta::info` for the purposes of argument-dependent lookup (ADL). That makes the use of various other facilities in that namespace considerably more convenient. (In this sense, `std::meta::info` is similar to an enumeration type.)

By requiring the type to be scalar, we avoid implementation overheads associated with the compile-time evaluation of class objects, indirection, and inheritance. By making the type unspecified but distinct, we avoid accidental conversions to other scalar types, and we gain the ability to define core language rules that deal specifically with these values. Moreover, no special header is required before using the lifting operator.

Reflection categories

As noted earlier, reflection values behave as handles to internal structures of the compiler. To reason about the kind of semantic information one can obtain through these reflection values, we categorize the values into one or more of four groups:

- Declared-entity reflections
- Alias reflections
- Expression reflections
- Invalid reflections

Note, declared-entity-reflections *only* designate the declared-entity; alias-reflections always designate a declared-entity in addition to providing the name of the alias; and, expression-reflections might or might not designate a declared-entity (e.g., an *id-expression* might designate a variable), but always designate properties of the expression. *Invalid reflections* will be discussed in more detail later, but they represent various kinds of failures when creating reflections using means other than the `^` operator.

For the most part, reflections of *names* (including *type-ids*) designate the declared entity those names denote: variables, functions, types, namespaces, templates, etc. For example:

```
^const int           // Designates the type const int.
^std                // Designates the namespace std.
^std::pair          // Designates the template pair.
^std::pair<int, int> // Designates the specialization.
int* f(int);        // Designates the type int*.
^decltype(f(3))     // Designates the type int*.
```

Reflections of *expressions* designate a limited set of characteristics of those expressions, including at least their type and value category. For example:

```
^1 // Designates the property “prvalue of type int” (but also the constant value 1)..
```

(Further on we will present facilities to examine and/or splice the designated notions.)

If an expression also names a declared entity (via a possibly-parenthesized *id-expression*), then it also designates that entity. For example:

```
int x;
^(x)      // Designates the declared-entity 'x' (variable) as well as the properties of
           // the expression 'x' (type and value category, in this case).
^(x+1)    // Does not designate a declared-entity but does designate the property
           // “prvalue of type int” (if ‘x+1’ had been constant-valued, it would also
           // designate the value it represents).
^std::cout // Designates the object named by std::cout as well as the
           // type and value category (lvalue) of the expression.
```

If an expression is a *constant expression* it also designates that constant value:

```
^0          // Designates the value zero and the property “prvalue of
           // type int”. It does not capture that the expression is a
           // is a literal or that it is usable as a null pointer value.
^nullptr   // Designates the null pointer value and the property “prvalue
           // of type decltype(nullptr)”.
^std::errc::bad_message // Designates the enumerator, its constant value, and the
           // property “prvalue of type std::errc”.
```

If an expression represents a call at its top level, it also designates the function being called (but not, e.g., the arguments to that call):

```
^printf("Hello, ")           // Designates printf and the property “prvalue of type
                           // int”.
^(std::cout << "World!")    // Designates the applicable operator<<
                           // and “lvalue of type std::ostream”.
constexpr int f(int p) { return p+1; }
^f(41)                     // Designates f, the (returned) value 42, and
                           // “prvalue of type int”.
^(f(41)+1)                 // Designates the (returned) value 43 and
                           // “prvalue of type int”; does not designate f
                           // because the call is not “top level”.
```

Now consider:

```
constexpr int const i = 42;
constexpr auto r = ^i;
```

As mentioned before, reflections can be categorized into four groups: declared-entity, alias, expression, or invalid. In this example, the reflection value `r` is an “expression reflection” and thus designates both the *expression* `i` (i.e. you can obtain information about properties of the expression such as its lvalueness) and the *variable* `i`. However, sometimes it is useful to obtain a reflection that designates only the entity (and not the expression). For example, we might want to query the type of the *variable* `i` (`int const`) instead of the type of the *expression* `i` (`int`). It also can be useful when comparing if two reflections refer to the same entity, as we will show later. We therefore provide the special function⁵

```
namespace std::meta {
    consteval auto entity(info reflection)->info {...};
}
```

which when applied to `r` produces a reflection designating just the *variable* (i.e., a “declared-entity reflection”).

More generally, `std::meta::entity` extracts the declared-entity from its argument by returning:

- its argument — if its argument is a declared-entity reflection or an invalid reflection,
- a declared-entity reflection designating an entity E — if the argument is an alias or expression reflection that also designates E, or
- an invalid reflection in all other cases (e.g., `entity(^42)` is an invalid reflection).

When the `^` operand is the name of an *alias* (type or namespace) the reflection designates the aliased entity indirectly (i.e., properties of the alias can be queried directly). For example:

```
using T0 = int;
using T1 = const T0;
constexpr meta::info ref = ^T1;
```

Here, `ref` designates both `T1` (directly) and the type `const int` (indirectly). This allows users to work both with the alias and its meaning. However, underlying aliases are not designated: There is no way to find about `T0` through `ref`.

In a more abstract sense, reflections designate semantic notions (names, types, value categories, etc.) rather than syntax (tokens that comprise an expression and the relation of those tokens to others). This principle helps guide decisions about the design of language and library support for reflection.

⁵ We use trailing return types for standard meta functions, but that’s just a stylistic preference. The traditional return type style is just as valid.

The queryable properties of these reflections are determined by the kind of “thing” they reflect. Details are provided below.

Equality and equivalence

Reflections can be compared using `==` and `!=` operators. Intuitively, the rule for these comparisons is that we compare the underlying declared entity, except that we cannot compare the reflections of most expressions nor can we compare invalid reflections. The exact rules are as follows...

1. If two reflections designate declared entities or aliases of such entities and do not designate expression properties of an expression that is not an *id-expression*, the reflections compare equal if the entities are identical and unequal if the entities are not identical (i.e., the comparison “looks through” aliases).
2. Any reflection also (obviously) compares equal to itself and to copies of itself.
3. An invalid reflection compares unequal to a reflection that is not invalid.
4. A reflection that designates a declared entity or an alias of such an entity and does not designate expression properties of an expression that is not an *id-expression* (e.g., it is not the reflection of a function call) compares unequal to a reflection that either does not designate an entity or an alias of such entity, or that designates properties of an expression that is not an *id-expression*.
5. All other cases are unspecified: That includes comparing reflections of expressions other than *id-expressions* and invalid reflections. For example:

```

typedef int I1;
typedef int I2;
static_assert(^I1 == ^I2);      // Rule 1: Same underlying declared entity (int).
static_assert(^I1 == ^int);    // Ditto.

float f = 3.0, e;
static_assert(^f == ^(f));    // Rule 1: Same underlying declared entity (f).
static_assert(^f == ^::f);    // Ditto.
static_assert(^f != ^e);     // Rule 1: Different underlying declared entities.
static_assert(^I1 != ^float); // Ditto.

void g(int);
constexpr auto r = ^g(1), s = r;
static_assert(r == s);        // Rule 2: One is a copy of the other.
static_assert(^f != ^g(1));   // Rule 4: f is an id-expression and g(1) is not.
static_assert(^g != ^g(1));   // Rule 4: One is the reflection of an id-expression
                           // and the other is the reflection of an expression
                           // that is not an id-expression.

static_assert(^g(1) == ^g(1)); // Rule 5: May fail because g(1) is an expression
                           // that is not an id-expression.

```

Programmers can more precisely specify whether they intend to compare entities or computed values (if possible) using splicers (e.g., `typename[:r:]` vs. just `[:r:]`) or library facilities like `std::meta::entity` described above. For example:

```
void f();    // #1
int f(int); // #2
constexpr auto r = std::meta::entity(^f(42)); // Designates function #2.
static_assert(r != ^f(42));                  // Fails.
static_assert(r == entity(^f(0)));           // Always succeeds..
```

Note that rule 1 above also applies to namespace aliases:

```
namespace N {};
namespace N1 = N;
namespace N2 = N;
static_assert(^N1 == ^N2);
static_assert(^N1 == ^N);
namespace M {};
static_assert(^N != ^M);
```

For reflections obtained from operands that involve template parameters, the result depends on the template arguments used for substitution:

```
template<typename T, typename U> struct Fun {
    static_assert(^T == ^U);
};
Fun<int, int> whee; // Okay.
Fun<int, char> oops; // Error: static assertion fails.
```

We already mentioned that it is unspecified whether reflections obtained from expressions that do not designate a declared entity compare equal. That also applies to expressions that just consist of a literal. For example:

```
static_assert(^1 == ^1); // May or may not fail.
```

(These rules allow us to avoid having to provide a general definition of “expression equivalence”.)

Note that the properties associated with a declared entity may change over various contexts, but that does not change the reflection. For example:

```

struct S;
constexpr auto r1 = ^S;
struct S {};
constexpr auto r2 = ^S;
static_assert(r1 == r2);

```

However, queries against the reflection value (e.g., to obtain a list of class members) may change as a consequence of the changes in the underlying entity.

An additional comparison function is proposed:

```

namespace std::meta {
    consteval auto same_reflections(info, info)->bool { ... };
}

```

If either `x` or `y` designate an alias (type or namespace) `same_reflections(x, y)` returns `true` if `x` and `y` designate the same alias and `false` otherwise. Otherwise (i.e., if neither `x` nor `y` designate an alias), `same_reflections(x, y)` returns `x == y`. In other words, `same_reflections(x, y)` is like the equality operator except that it doesn't "look through" aliases. For example, with the namespace aliases `N1` and `N2` as above:

```

using std::meta::same_reflections;
static_assert(!same_reflections(^N1, ^N2));
static_assert(same_reflections(^(^N1), ^(^N1))); // May fail.
static_assert(!same_reflections(^(^N1), ^(^N2))); // May fail.

```

The latter two assertions have unspecified behavior because `^N1` (or `^N2`) is an expression that is not an *id-expression* and (as was noted above) the equality of the reflections of such expressions is unspecified.

To compare the values of reflected objects, references, functions, or types, the reflection can first be spliced (see below).

A Note About Linkage

Although in most respects we propose that `std::meta::info` is an ordinary scalar type, we also give it one "magical" property with respect to linkage.

Before explaining this property, consider again what a reflection value represents in practice: It is a handle to internal structures the compiler builds up for the current translation unit. So for code like:

```

struct S {};
consteval auto f() {
    return ^S;
}

```

the compiler will construct an internal representation for struct **S** and when it encounters “**^S**” it will update a two-way map between the internal representation of **S** and a small structure underlying the **std::meta::info** value returned by **^S**.

Now consider:

```

// Header t.hpp:
struct S {};
template<std::meta::info reflection> struct X {};

// File t1.cpp:
#include "t.hpp"
enum E {};
consteval auto d() {
    return ^E;
}
X<^S> g() {
    return X<^S>{};
}

// File t2.cpp:
#include "t.hpp"
extern X<^S> g();
int main() {
    g();
}

```

The files **t1.cpp** and **t2.cpp** are compiled separately. The contexts in which the “**^S**” construct is encountered are therefore different and it is not practical to ensure that the *underlying* values (“bits”) of the **std::meta::info** results are identical. However, it is *very* desirable that the types **X<^S>** are the same types in both translation units and that the above example *not* produce an ODR violation.

We therefore specify “by fiat” that:

- **reinterpret_cast** to or from **std::meta::info** is ill-formed
- accessing the byte representation of **std::meta::info** lvalues produces unspecified (possibly inconsistent) values
- **std::meta::info** values **A1** and **A2** produce equivalent template arguments if **std::meta::same_reflections(A1, A2)** produces **true**.

However, it is unspecified if the following variation of the previous example is valid:

```
// File t1.cpp:  
enum E {};  
consteval auto d() {  
    return ^(^E);  
}  
X<^(^S)> g() {  
    return X<^(^S)>{};  
}  
  
// File t2.cpp:  
extern X<^(^S)> g();  
int main() {  
    g();  
}
```

because it is unspecified if two occurrences of `^(^S)` are equivalent.

(In practice, this means that reflection values are mangled symbolically, according to what the reflection value actually designates.)

Invalid reflections

In what follows we are going to propose a large collection of standard reflection operations, some of which generate new reflection values. Sometimes, the application of some of these operations will be meaningless. E.g., consider:

```
namespace std::meta {  
    consteval auto add_const(info)->info {...};  
}
```

which is meant to take a reflection of a type and add a type qualifier on top. However, what happens with something like:

```
constexpr auto r = add_const(^std);
```

which suggests the meaningless operation of adding a `const` qualifier to namespace `std`? Our answer is that an implementation will not immediately trigger an error in that case, but instead create a reflection value that represents an error. Any attempt to splice such a reflection is ill-formed (but subject to SFINAE).

It is useful for user code to also be able to produce invalid reflections. To that end, we propose the

following function:

```
namespace std::meta {
    consteval
        auto invalid_reflection(std::string_view message,
                                std::source_location src_loc =
                                std::source_location::current())
            ->info {...};
}
```

which constructs a reflection that triggers a diagnostic if it is spliced outside a SFINAE context (ideally, with the given message and source location information). Here, the functions

Invalid reflections can also be used to generate compiler diagnostics during constant evaluation using the `diagnose_error` function. This can be a valuable debugging aid for authors of metaprogramming libraries, and when used effectively, should improve the usability of those libraries.

```
namespace std::meta {
    consteval void diagnose_error(info invalid_refl) {...};
}
```

This function causes the compiler to emit an error diagnostic (formally: it makes the program ill-formed if it is invoked outside a deduction/SFINAE context), hopefully with the message and location provided by the argument. For example:

```
auto r = std::meta::invalid_reflection("Oops!");
int main() {
    diagnose_error(r); // Error.
}
```

That last example is ill-formed and might trigger an error like:

```
"Test.cpp", line 3: error: Invalid reflection
diagnose_error(r); // Error.
^
"Test.cpp", line 1: note: Oops!
auto r = std::meta::invalid_reflection("Oops!");
^
```

Finally, we propose a predicate:

```
namespace std::meta {
    consteval auto is_invalid(info)->bool {...};
}
```

that can be used to, e.g., filter out invalid reflective operations. We also provide a convenience overload of this function:

```
namespace std::meta {
    consteval auto is_invalid(span<info>)->bool {...};
}
```

which returns `true` if any element of the given span is an invalid reflection. This is particularly useful because some important reflection facilities return spans of reflection values that callers are likely to want to check for invalid entries.

Initialization of reflections

Objects of reflection type are zero-initialized to an invalid reflection value (with unspecified associated information).

Conversions on reflections

A prvalue of reflection type can be contextually converted to a prvalue of type `bool`. An invalid reflection converts to `false`; all other reflections convert to `true`.

Hashing reflections

We propose that the `std::hash` template be specialized for `std::meta::info`. We also propose that the resulting hash value be consistent across translation units.

Splicing

In the context of this paper, “splicing” refers to the process of turning a “reflection value” back into a “program source thing”. We propose a basic *splice* construct to be of the form

`[: reflection :]`

where `[:` and `:]` are each a sequence of two tokens and `reflection` is a constant-expression of type `std::meta::info`. (Prior versions of this paper discuss various alternative syntax options. The choice presented here was first proposed in P2320R0, which obtained strong support in SG-7.)

In general, and without qualification, `[: R :]` splices an expression into the program (assuming `R` reflects a variable, function, or a constant expression). If `R` reflects both a constant value and a declared entity

(e.g., `^f()` where `f` is a `consteval` function), `[: R :]` splices the constant value. Use `std::meta::entity` to get the other outcome. If `R` reflects a type or template, the splice construct must be qualified with an appropriate `typename` or `template` keyword, except in some contexts where the meaning is obvious. For example:

```
struct S { struct I { }; };
template<int N> struct X;
auto refl = ^S;
auto tmpl = ^X;
void f() {
    typename[ :refl:] * x;      // Okay: declares x to be a pointer to S.
    [ :refl:] * x;            // Error: attempt to multiply int by x.
    [ :refl:]::I i;          // Okay: splice as part of a nested-name-specifier.
    typename[ :refl:]{};      // Okay: default-constructs an S temporary.
    using T = [ :refl:];      // Okay: operand must be a type.
    struct C: [ :refl:] {};   // Okay: base classes are types.
    template[ :tmpl:]<0>;    // Okay: names the specialization.
    [ :tmpl:] < 0 > x;       // Error: attempt to compare X with 0.
}
```

When a splice construct is used as a template argument, without a disambiguating `typename` or `template` keyword, is not *a priori* assumed to be an expression. For example:

```
template<typename T> void f();
template<int N> void f();
template<meta::info Refl> void ex1() {
    f<[ :Refl:]>(); // Could resolve to either function template above.
}
```

To “force” `Refl` to be spliced as a type or an expression by adding `typename` or enclosing the splice in parentheses:

```
template<meta::info Refl> void ex2() {
    f<typename [ :Refl:]>(); // splices a type
    f<([ :Refl:])>(); // splices an expression
}
```

Here are some additional examples illustrating various uses of the splicing constructs:

```
typename[ :^int:] i = [ :^42:]; // Same as "int i = 42;".
```

```
constexpr int J = 42;
[ :^i:] = [ :^J:]; // Same as "i = J;".
```

```

[ :^i:] = [ :^(J+1):]; // Same as "i = 43;".

[ :^i:] = [ :^[]{ return J; }():];
// Same as "i = 42;".

[ :^i:] = [ :^ [=]{ int x = i; return J; }():];
// Error: The lambda call is not a constant expression.
// (Reflections of arbitrary expressions cannot be spliced.)

namespace N { int f; }

void [ : ^N::f :](int);
// Error: Not an expression context.

struct S {
    consteval auto ri() { return ^S::i; }
private:
    int i:3; // Bit field.
} s;
int i1 = s.[ :s.ri():]; // Okay: Refers to S::i without needing name
                        //           lookup at this point.

```

Furthermore, we propose a sequence-generating splicing construct. Let *reflection_range* be a *constant range* such that

```
for (std::meta::info r : reflection_range) ...
```

would successively set *r* to a sequence of values *r*₁, *r*₂, *r*₃, ... *r*_N.

Then

```
... [ : reflection_range :] ...
```

expands like a parameter pack, but unlike a parameter pack it might be a heterogeneous expansion in that some elements of the expansions might be a type, a constant expression, or a template. The prefix ... token turns the subsequent splicer into a “pack-like” construct, and the trailing ... does the expansion as usual.

Examples:

```

std::meta::info t_args[] = { ^int, ^42 };
template<typename T, T> struct X {};
X<...[:t_args:]...> x; // Same as "X<int, 42> x;".
template<typename, typename> struct Y {};
Y<...[:t_args:]...> y; // Error: same as "Y<int, 42> y;".

```

Some observations:

- Empty ranges and singleton ranges expand as expected.
- If any expansion produces an ill-formed splice, the whole construct is ill-formed but subject to SFINAE.

Splicing a function-local alias or declared entity outside its potential scope is ill-formed. For example:

```
consteval auto refl_int_alias() {
    typedef int Int;
    return ^Int;
}
typename[:refl_int_alias():] x; // Error: Cannot splice local alias here.
```

Similarly, a parameter obtained from a function type F can be spliced as an expression only within the potential scope of the corresponding argument of a function of the same type. For example (`parameters_of` will be described later on):

```
using F = int (int, int);
auto params = std::meta::parameters_of(^F);
int f(int p, int) {
    return [:params[0]:]; // Okay: Same as "return p;"
}
int g(int, char) {
    return [:params[0]:]; // Error: params[0] comes from function type
} // "int (int, int)" but this function has
   // type "int (int, char)".
```

Splicing identifiers

We anticipate the later addition of an *identifier-splice* construct (currently we use the `[# str #]` syntax in discussions among authors). However, that construct operates, in part, at the lexical level and has considerably more subtleties than the authors are exploring (in part through prototype implementations). We therefore do not propose syntax for it here, and we expect that the corresponding functionality will be proposed separately later on. Unlike prior versions of this paper, this revision does not yet propose the addition of such a capability.

Access checking

Splicers provide an alternative way to refer to declarations and therefore we must decide whether they are subject to access control. Access control ordinarily applies to *names*, but the `[:...:]` construct does not create a name and is thus *not* subject to access checking. For example:

```

class C {
    using Int = int;
public:
    static consteval auto r() { return ^Int; };
} c;
typename[ :C::r():] x; // Okay: x has type int

```

If we introduce *identifier splicing* later on (as suggested above), those kinds of splicers *would* be subject to access checking.

Templates and reflection

Reflection mostly occurs “after instantiation”. For example, in:

```

template<typename T> void f(T p) {
    constexpr auto r = ^T;
}

```

the expression T is always a *dependent* expression that doesn’t produce an actual value until f is instantiated. I.e., this does not provide a mechanism to get a handle on a reflection for the template parameter T itself. However, that doesn’t mean that we don’t propose *any* facilities to reflect templated entities.

Template arguments

We propose a function

```

namespace std::meta {
    consteval
    auto has_template_arguments(info reflection)->bool {...};
}

```

that returns `true` if and only if the given reflection corresponds to a template specialization (in the standard sense: implicit specializations are included).

The actual template arguments can be obtained through

```

namespace std::meta {
    consteval auto template_arguments_of(info reflection)
        ->std::span<info> {...};
}

```

Conversely, the template producing a specialization can be obtained with

```
namespace std::meta {
    consteval auto template_of(info reflection)->info {...};
}
```

Note that the resulting reflection value (like that for reflecting a template directly) represents that template as completely known at any point it is examined (including not only the primary template definition, but also partial and full specializations). If the given reflection is not that of a specialization, an invalid reflection is returned.

Template substitution

We also propose a facility that is the “dual notion” of the previous functions:

```
namespace std::meta {
    consteval auto substitute(info templ, std::span<info> args)
        ->info { ... };
}
```

This produces a declared-entity reflection for an instance given a reflection for a template and a span of reflections for specific arguments. A substitution error in the immediate context of the substitution produces an invalid reflection (this is akin to SFINAE). A substitution error outside that immediate context renders the program ill-formed. An incomplete substitution (where not all parameters are substituted by nondependent arguments) also produces an invalid reflection. Note, this functionality can also be approximated using splicers with

```
^template[ :templ:]<...[ :args:]...>
```

but having both improves readability depending on the context. The **substitute** form has the added advantage of not triggering an error for failures in the immediate context of the substitution⁶.

Example:

```
using namespace std::meta;
template<typename ... Ts> struct X {};
template<> struct X<int, int> {};
constexpr info type = ^X<int, int, float>;
constexpr info templ = template_of(type);
constexpr span<info> args = template_arguments_of(type);
constexpr info new_type = substitute(templ, args.subspan(0, 2));
```

⁶ While working with our implementations, we have noticed that it would be very convenient if the lifting operator would be a SFINAE context as well. E.g., instantiating `^T:X` would produce an invalid reflection when `T = int`. That option is still being considered.

```
typename[ :new_type:] xii; // Type X<int, int>, which selects the specialization.
                           // There is no mechanism to instantiate a primary template
                           // definition that is superseded by an explicit/partial
                           // specialization.
```

The use of substitute in that last example could instead be written as:

```
constexpr info new_type =
    ^template[ :templ:]<...[ :args.subspan(0, 2):]...>;
```

Another example illustrates how substitutions could produce non-SFINAE errors:

```
template<typename T> struct A {
    T::type I;
};

template<typename T, T::type N> struct Y {};
constexpr info ASpec = ^A<int>; // No instantiation yet.
constexpr info new_type2 =
    substitute(^Y, std::vector<info>{ ASpec, ^5 });
                           // Error: Substitution of Y<A<int>, 5> requires A<int> to be instantiated
                           //           outside the immediate context of the substitution.
```

Template parameters

Although applying `^` to dependent constructs doesn't produce an actual value until instantiation/substitution of the enclosing templated entity, reflections of template parameters can be obtained from the reflection of a template:

```
namespace std::meta {
    consteval auto parameters_of(info reflection)
        ->std::span<info> {...};
}
```

Given the reflection of a template, this returns a sequence of reflections for each template parameter. Each of these reflections can be a type, a constant, or a template. However, not all operations applicable to types/constants/templates are necessarily applicable to these reflections. For example, it would not be possible to apply the `std::meta::substitute` operation (when available) on the reflection of template template parameters (but it *is* possible to apply the `std::meta::parameters_of` to such a reflection).

The standard metaprogramming library

We have already described a number of “metafunctions” living in namespace `std::meta` that work with reflections (`entity`, `invalid_reflection`, `parameters_of`, etc.). We propose that these and many others make up a new section of the C++ standard library that we call “the C++ standard metaprogramming library” and which is made available to a program by including a new standard header `<meta>`.

Besides the metafunctions that we have described so far, we add additional functions based on two sources:

- the existing standard library’s [meta] section (described in the “General utilities” clause of the current working paper), and
- the [reflect] section of the original “Reflection TS” (N4856).

We delve into how to “transcribe” those additional metafunctions from the original source below, but first a note about the API design.

Sequence and string values

Many metafunctions deal with sequences (particularly, sequences of reflection values) and with string values (particularly, names of entities). Earlier versions of this paper used `std::vector` and `std::string` types for this. However, that was counting on those types being available for constant evaluation. P0784R7 mostly achieved that availability, but the inability to have nontransient dynamic allocation means that, e.g., using `std::string` would make it difficult to transfer a reflection string into the run-time domain.

We therefore now use `std::span` and `std::string_view` types for this, since they do not require dynamic allocation of backing storage. Instead, the compiler will have to provide static storage for, e.g., reflected strings. In practice, an implementation will want to be careful to only emit those strings in object code if they are actually used in the run-time domain, but that is not a particularly difficult implementation challenge.

Transcribing the standard library’s [meta] section

The standard library [meta] section (in clause [utilities]) provides a large number of utilities to examine and construct types. We propose that all those utilities be given a counterpart in the value-based reflection world, with needed declarations made available through a new standard header `<meta>`.

For example, consider the type transformation trait

```
std::make_signed<T>
```

which produces a result through its member type

```
std::make_signed<T>::type
```

We propose to have a `std::meta` counterpart as follows:

```
namespace std::meta {
    consteval auto make_signed(info reflection)->info {...};
}
```

This is expected to be implemented using an intrinsic in the compiler (although that is not a requirement). For a reflection value `r` corresponding to a type `T` such that

```
std::make_signed<T>::type
```

is valid, using the new function as `make_signed(r)` is equivalent to:

```
^std::make_signed<typename(r)>::type
```

(except for not actually instantiating templates in a quality implementation). For a reflection value for which the above transformation would not be valid (e.g., `^void`), however, the function returns an invalid reflection.

Most templates specified in [meta.trans] can be transcribed in a similar way, but a few take additional nontype template parameters. Their transcription is also straightforward however. We illustrate this with the `std::enable_if` template whose `consteval` counterpart can be implemented efficiently without intrinsics. The already-standard template-based interface is usually implemented as follows:

```
namespace std {
    template<bool, typename T = void> struct enable_if {};
    template<typename T> struct enable_if<true, T> {
        using type = T;
    };
}
```

The reflection counterpart is then (including a hypothetical implementation):

```

namespace std::meta {
    consteval auto enable_if(bool cond, info type = ^void)->info {
        if (cond) {
            return type;
        } else {
            return invalid_reflection("enable_if condition false");
        }
    };
}

```

(We encourage programmers to prefer *requires-clauses* over `enable_if` for constraining templates.)

The type traits predicates described in [meta.unary] and [meta.rel] are just as easily mapped to the value-based reflection world. For example, the three templates

```

namespace std {
    template<typename T> struct is_union;
    template<typename T, typename ... Args> struct is_constructible;
    template<typename B, typename D> struct is_base_of;
}

```

have counterparts as follows:

```

namespace std::meta {
    consteval auto is_union(info reflection)->bool {...};
    consteval auto is_constructible(info reflection,
                                   std::span<info> arg_types)
                           ->bool {...};
    consteval auto is_base_of(info base_type,
                            info derived_type)->bool {...};
}

```

The other cases follow the same patterns.

The three templates in [meta.unary.prop.query]:

```

namespace std {
    template<typename T> struct alignment_of;
    template<typename T> struct rank;
    template<typename T, unsigned I = 0> struct extent;
}

```

are slightly irregular, but the corresponding functions can still be intuited:

```
namespace std::meta {
    consteval auto alignment_of(info type)->std::size_t {...};
    consteval auto rankinfo(type)->int {...};
    consteval auto extent(info type, unsigned dim = 0)->int {...};
}
```

The helper templates in [meta.help] and [meta.logical] are not needed for value-based reflection since their counterparts are core language features (like the integer types and the logical operators).

Adapting the Reflection TS' [reflect] section

The Reflection TS (N4818) introduces a large number of template metafunctions. This proposal steals many of those features and adapts them to the value-based reflection world. However, we make some changes to better align the semantics with the constraints of the language definition and the flexibility of our value-based approach.

Predicates

Let's start with the predicates (metafunctions returning a `bool` value). For example, `is_public` gets a counterpart as follows:

```
namespace std::meta {
    consteval auto is_public(info base_or_mem)->bool {...};
}
```

That function fails to evaluate to a constant if `base_or_mem` does not designate a base class or a class member (that constraint corresponds to the concepts requirements imposed for the class template `is_public` proposed in the Reflection TS). `is_protected`, `is_private`, `is_accessible` (which checks whether a member is accessible from the context of invocation), `is_virtual`, and `is_final` are handled in the same way. For example:

```
struct S { int x; };
constexpr bool t = std::meta::is_public(^S::x);
                // = true;
constexpr bool e = std::meta::is_public(^S);
                // Error: Not a constant because ^S is not a base or member.
```

The `is_unnamed` metafunction is transcribed similarly:

```
namespace std::meta {
    consteval auto is_unnamed(info entity)->bool {...};
}
```

but this time the function only evaluates to a constant if the given reflection represents a namespace, a data member, a function, a template, a variable, a type, or an enumerator. Failing to evaluate a constant will result in an error, or a deduction failure if in a SFINAE context.

`is_scoped_enum` becomes

```
namespace std::meta {
    consteval auto is_scoped_enum(info entity)->bool {...};
}
```

and always evaluates to a constant.

We propose to replace `is_constexpr` by:

```
namespace std::meta {
    consteval auto is_declared_constexpr(info entity)->bool {...};
}
```

which is a constant value if `entity` designates a variable, a function, a static data member, or a template for these. (We propose the alternative name to distinguish the entities that are declared with the `constexpr` or `consteval` specifier from entities that are effectively `constexpr` (e.g., a function template may be declared `constexpr` and its instances would produce `true` values with this predicate; however, the instances may not actually be `constexpr` functions; conversely, lambda call operators and special member functions may be `constexpr` functions without being declared `constexpr`). Immediate (`consteval`) functions and function templates are also identifiable:

```
namespace std::meta {
    consteval auto is_consteval(info entity)->bool {...};
}
```

Instead of `is_static` (for variables) we propose:

```
namespace std::meta {
    consteval
    auto has_static_storage_duration(info entity)->bool {...};
}
```

because “`is_static`” suggests a query about a storage class specifier rather than a storage duration.

```
namespace std::meta {
    consteval auto is_inline(info entity)->bool {...};
}
```

produces a constant value for reflections of variables, functions, variable/function templates, and namespaces.

A number of function properties produce a constant value for reflections of functions only:

```
namespace std::meta {
    consteval auto is_deleted(info entity)->bool {...};
    consteval auto is_defaulted(info entity)->bool {...};
    consteval auto is_explicit(info entity)->bool {...};
    consteval auto is_override(info entity)->bool {...};
    consteval auto is_pure_virtual(info entity)->bool {...};
}
```

The following predicates always produce a constant value given a reflection. They produce a **false** value for invalid reflections, and otherwise return **true** if the predicate applies to the reflected entity:

```
namespace std::meta {
    consteval auto is_class_member(info reflection)->bool {
        // Return true for class and class template members.
        ...
    };
    consteval auto is_local(info reflection)->bool {
        // Return true for local variables, local members.
        ...
    };
    consteval auto is_namespace(info entity)->bool {...};
    consteval auto is_template(info entity)->bool {...};
    consteval auto is_type(info entity)->bool {
        // Return true for types and type aliases.
        ...
    };

    consteval auto is_incomplete_type(info entity)->bool;
    consteval auto is_closure_type(info entity)->bool {...};
    consteval auto has_captures(info entity)->bool {...};
    consteval auto has_default_ref_capture(info entity)->bool {
        // Return true even if there is no effective capture (i.e., it's syntactical only).
        ...
    };
    consteval auto has_default_copy_capture(info entity)->bool {
        // Return true even if there is no effective capture (i.e., it's syntactical only).
        ...
    };
}
```

```
consteval auto is_simple_capture(info entity)->bool {...};
consteval auto is_ref_capture(info entity)->bool {...};
consteval auto is_copy_capture(info entity)->bool {...};
consteval auto is_explicit_capture(info entity)->bool {...};
consteval auto is_init_capture(info entity)->bool {...};
consteval auto is_function_parameter(info entity)->bool {...};
consteval auto is_template_parameter(info entity)->bool {...};
consteval auto is_class_template(info entity)->bool {...};
consteval auto is_alias(info reflection)->bool {...};
consteval auto is_alias_template(info reflection)->bool {...};
consteval auto is_enumerator(info entity)->bool {...};
consteval auto is_variable(info entity)->bool {...};
consteval auto is_variable_template(info entity)->bool {...};
consteval auto is_static_data_member(info entity)->bool {
    return is_variable(entity) && is_class_member(entity);
};
consteval auto is_nonstatic_data_member(info entity)->bool {
    // Return true for nonstatic data members, which includes bit fields.
    ...
};

consteval auto is_bit_field(info reflection)->bool {
    // Return true for bit fields, but also for expressions that are bit field selections.
    ...
};

consteval auto is_base_class(info entity)->bool {...};
consteval auto is_direct_base_class(info entity)->bool {...};
consteval auto is_virtual_base_class(info entity)->bool {
    return is_base_class(entity) && is_virtual(entity);
};

consteval auto is_function(info entity)->bool {...};
consteval auto is_function_template(info entity)->bool {...};

consteval auto is_member_function(info entity)->bool {
    return is_function(entity) && is_class_member(entity);
};

consteval
auto is_member_function_template(info entity)->bool {
    return is_function_template(entity) && is_class_member(entity);
};

consteval
auto is_static_member_function(info entity)->bool {...};
consteval
auto is_static_member_function_template(info entity)->bool {...};
```

```

consteval
auto is_nonstatic_member_function(info entity)->bool {...};
consteval
auto is_nonstatic_member_function_template(info entity)
    ->bool {...};
consteval auto is_constructor(info entity)->bool {...};
consteval auto is_constructor_template(info entity)->bool {...};
consteval auto is_destructor(info entity)->bool {...};
consteval auto is_destructor_template(info entity)->bool {...};
} // namespace std::meta

```

Note that `is_bit_field` above is more general than what the TS proposed since it applies not only to the reflection of data members but also to expressions, because “bitfieldness” is a significant property of an expression. Similarly, we add the following five predicates (with no equivalent in the TS) for reflections of expressions:

```

consteval auto is_lvalue(info reflection)->bool;
consteval auto is_xvalue(info reflection)->bool;
consteval auto is_prvalue(info reflection)->bool;
consteval auto is_glvalue(info reflection)->bool {
    return is_lvalue(reflection) || is_xvalue(reflection);
}
consteval auto is_rvalue(info reflection)->bool {
    return is_prvalue(reflection) || is_xvalue(reflection);
}

```

The following predicate produces a constant value given the reflection of a function type or closure type, or an alias thereof:

```

namespace std::meta {
    consteval auto has_ellipsis(info entity)->bool {...};
}

```

The following predicate produces a constant value given the reflection of a function type or an alias thereof:

```

namespace std::meta {
    consteval auto is_member_function_type(info entity)->bool {...};
}

```

Given the reflection of a function or template parameter, `std::meta::has_default` returns whether it has an associated default argument:

```
namespace std::meta {
    consteval auto has_default(info entity)->bool {...};
}
```

Singular properties

This section lists facilities that return properties described with a “single value”.

The following function can be used to identify a source location of a declared entity:

```
namespace std::meta {
    consteval auto source_location_of(info entity)
                    ->std::source_location {...};
}
```

Although this produces a constant result for any reflection value, the returned value is unspecified if the reflection is not that of a declared entity (or alias).

The name of declared entities can be accessed through the following:

```
namespace std::meta {
    consteval auto name_of(info entity)->std::string_view {...};
    consteval auto display_name_of(info entity)
                    ->std::string_view {...};
}
```

For named declared entities/aliases, `name_of` returns a constant `string_view` describing the same identifier as that produced by the “[: info :]” splicer. For any other operand, it produces a constant empty `string_view`.

The `display_name_of` function produces an unspecified constant non-empty `string_view` for any reflection (implementations are encouraged to produce a string that is helpful in identifying the reflected item).

Aliases can be “looked through” using the aforementioned function `entity`:

```
namespace std::meta {
    consteval auto entity(info reflection)->info {...};
}
```

A reflection for the type associated with an entity or expression can be retrieved with

```
namespace std::meta {
```

```
    consteval auto type_of(info reflection)->info {...};
}
```

If reflection describes an entity (not an expression) that is not a variable, base class, data member, function, or enumerator, this function returns an invalid reflection.

A “parent” entity can be identified with

```
namespace std::meta {
    consteval auto parent_of(info reflection)->info {...};
}
```

For members of classes or namespaces this returns a reflection of the innermost class or namespace. For a base class, this returns the class type from which the base class was obtained (only direct and virtual base classes can be reflected). For function-local entities that are not class members, `parent_of` returns the a reflection of the enclosing function. For reflections that do not designate an alias or a declared entity, `parent_of` returns an invalid reflection.

The innermost enclosing function and class can also be queried:

```
namespace std::meta {
    consteval auto current_function()->info {...}
    consteval auto current_class_type()->info {...}
}
```

That is particularly useful to deal more efficiently with parameter packs (an example will be presented later on). Note that when invoked from an immediate function in a context that does not require a constant-expression, these functions return the result as if invoked from the calling function.

Given the reflection of a base or nonstatic data member of a class (but not a class template), layout information can be retrieved with the following functions:

```
namespace std::meta {
    consteval auto byte_offset_of(info entity)->std::size_t {...};
    consteval auto bit_offset_of(info entity)->std::size_t {...};
    consteval auto byte_size_of(info entity)->std::size_t {...};
    consteval auto bit_size_of(info entity)->std::size_t {...};
}
```

For reflections that do not designate a base or a nonstatic data member, this does not successfully produce a constant value. `byte_offset_of` returns the byte offset of the given base or nonstatic data member (within the parent class). For bit-fields, the offset of the first byte containing the bit field is returned; the bit offset of the first bit (counting from the least significant bit) within that byte is produced by

`bit_offset_of` (for non-bit-fields, that function returns zero). `byte_size_of` produces the allocated size of the associated subobject, except that it does not produce a constant value for bit fields (for base classes, the result may be less than `sizeof` applied to the base class type). `bit_size_of` produces the allocated size of the associated bit field subobject, and does not produce a constant value for non-bit-field reflections. (A precise specification of this requires a slight tightening of the C++ object model. All implementations already conform to the stricter model.)

The following facilities permit examining parameter types and the “`this`” binding type:

```
namespace std::meta {
    consteval auto this_ref_type(info func_type)->info {...};
}
```

For a member function type `this_ref_type` returns the reflection of the parent class associated with the member type, with any member function *cv-qualifiers* and *ref-qualifiers* added on top. For example:

```
struct S {
    int f() volatile &&;
    int g() const;
} s;
constexpr auto r = this_ref_type(^s.f());
// Reflection for type "S volatile &&".
constexpr auto r = this_ref_type(^s.g());
// Reflection for type "S const".
```

Plural properties

This section lists metafunctions that return `std::span<info>` values describing “plural properties” (such as lists of members).

We propose the following function templates to retrieve subobject information:

```
namespace std::meta {
    template<typename ...Fs>
    consteval auto members_of(info class_type, Fs ...filters)
        ->std::span<info> {...};

    template<typename ...Fs>
    consteval auto bases_of(info class_type, Fs ...filters)
        ->std::span<info> {...};
}
```

If called with an argument for `class_type` that is the reflection of a non-class type or a capturing closure type (or an alias/cv-qualified version thereof), these facilities return a span referencing a single

invalid reflection.

Otherwise, if no `filters` argument is passed to it, `members_of` returns an “unfiltered” sequence of reflections for the following kinds of direct members of a class type (represented by `class_type`): nonstatic and static data members and member functions, member types (enumeration and class types) and member aliases, and member templates other than deduction guides. Generated members are included, but inherited constructors, injected-class-names, and unnamed bit fields are not (the standard doesn’t consider those members either). Nonstatic data members appear in declaration order (but not necessarily consecutively).

If any “filters” are passed, they are applied as predicates to the unfiltered sequence, and members for which a predicate produces `false` are left out. Predicates are applied left-to-right with short-circuit semantics (i.e., later predicates are not applied if an earlier predicate produced `false`).

Similarly, without `filter` arguments, invoking `bases_of` returns a sequence of reflections for the direct bases of the given (non-capturing-closure) class type. Predicates can be added to narrow down the bases of interest.

The following example illustrates some uses of `members_of`:

```
struct S {
    double x;
    int y;
    void f();
};

constexpr auto class_type = ^S;
constexpr auto s_members = members_of(class_type);
static_assert(s_members.size() == 7);
    // x, y, f(), the destructor, and generated constructors.

constexpr auto s_data_members =
    members_of(class_type, is_nonstatic_data_member);
static_assert(s_data_members.size() == 2);
    // x and y.

constexpr auto has_integral_type(std::meta::info reflection) {
    return std::meta::is_integral(std::meta::type_of(reflection));
}

constexpr auto s_imembers =
    members_of(class_type, is_nonstatic_data_member,
               has_integral_type);
static_assert(s_imembers.size() == 1);
    // Just y.

constexpr auto s_nested_types =
```

```

        members_of(class_type, is_type);
static_assert(s_members.size() == 0);
// S has no nested types.

```

The enumerators of an enumeration type can be inspected using `enumerators_of`:

```

namespace std::meta {
    consteval
    auto enumerators_of(info enum_type)->std::span<info> {...};
}

```

If the argument passed for `enum_type` is not a reflection for an enumeration type, this returns a span referencing just an invalid reflection.

The parameters of a function type or the parameters of a template can be inspected using:

```

namespace std::meta {
    consteval auto parameters_of(info reflection)
        ->std::span<info> {...};
}

```

If the argument passed for `reflection` is not a reflection for a function, a member function, a function type, a closure type, or a template, this returns a span referencing one invalid reflection. Otherwise, the span contains an entry for each ordinary parameter: No entry is made for the “`this`” parameter or for an ellipsis parameter.

A function is also available to introspect lambda captures associated with a closure type:

```

namespace std::meta {
    consteval auto captures_of(info closure_type)
        ->std::span<info> {...};
}

```

If the argument passed for `closure_type` is not a reflection for a closure type, this returns a span referencing just an invalid reflection.

Of note here is that we are *not* proposing a function to retrieve members of namespaces: Due to their “open scope” nature, we believe that capability is somewhat meaningless. There is however no known technical reason preventing us from doing so.

We are also *not* proposing functions to expose the structure of templates. In particular, there is no mechanism to retrieve the members of a class template or the function parameters of a function template.

Anonymous unions

Consider:

```
struct S {
    bool flag;
    union {
        int x;
        float f;
    };
};

constexpr auto dmembers =
    members_of(^S, is_nonstatic_data_member);
static_assert(dmembers.size() == 2);
```

The sequence `dmembers` here will contain two reflections: One for `flag` and one for an unnamed data member of the unnamed union type. Conversely, `parent_of(^S::x)` produces a reflection for that unnamed union type rather than for `S`. Despite its lack of a declared name (which means `name_of` returns an empty `string_view`), the unnamed data member can be referred to with the “[: ... :]” splicer:

```
constexpr S s = { false, { .x = 42 } };
static_assert(name_of(dmembers[1]) == ""); // Okay.
static_assert(s.idexpr( dmembers[1] ).x == 42); // Okay.
static_assert(
    remove_reference(^decltype(s.[ :dmembers[1]:])) ==
        parent_of(^S::f)); // Okay.
```

Let's take that last line apart.

In the left-hand side of the equality test `dmembers[1]` is a reflection of the unnamed data member for the anonymous union. Therefore, `s.[:dmembers[1]:]` is an lvalue designating the anonymous union subobject of `s`, and thus `decltype` applied to that produces a reference to the anonymous union type. The `^` operator returns the reflection designating that type and `remove_reference` finally returns a reflection designating the underlying union type.

In the right-hand side, `^S::f` is a reflection designating the member `S::f`, which is actually a member `S::<unnamed-union-type>::f`. Therefore, `parent_of` also produces a reflection designating the underlying union type of the anonymous union, and the assertion succeeds.

Other Facilities

Reflecting values

It turns out to be useful to be able to lift a constant value into reflections for an expression denoting that constant value. We therefore propose a pair of metafunction templates to do exactly that:

```
namespace std::meta {
    template<typename T>
        consteval auto reflect_value(T const&)->info;
    template<typename R>
        consteval auto reflect_values(R const&)->std::span<info>;
}
```

(The second template applies the first to each element of a range of values.)

For example:

```
constexpr std::vector<int> v{ 1, 2, 3 };
constexpr std::span<std::meta::info> rv = reflect_values(v);
```

Such a lifted sequence can then be spliced into a template argument context:

```
std::integer_sequence<int, ...[ :rv:]...> is123;
// same as std::integer_sequence<int, 1, 2, 3>
```

provided the reflected constants are valid in that context.

Metaprogramming Examples

We believe that the facilities presented here permit the kind of computation previously performed with C++ template metaprogramming and that they are preferable over TMP because they scale better. We therefore suggest that no broad set of TMP facilities should be further added to the language.

Examples in this section are drawn from a variety of sources, including P0385R0 by Matúš Chochlík and Axel Naumann and P0949R0 by Peter Dimov.

Hashing

We can also use the approach above to synthesize an overload of `hash_append` (proposed by Howard Hinnant *et al.* in [N3980](#), *Types Don't Know #*).

```
#include <meta>
```

```

using namespace meta = std::meta;
template<HashAlgorithm H, StandardLayoutType T>
bool hash_append(H &algo, const T &t) {
    constexpr auto data_members =
        members_of(^T, meta::is_nonstatic_data_member);
    template for (constexpr meta::info member : data_members)
        hash_append(algo, t.[:member:]);
}

```

The algorithm is straightforward: Recursively apply `hash_append` to each member for the class `T`. Within that call the expression `t . [:member :]` yields a *postfix-expression* for the designated member in the class object. The resolution of that postfix-expression does not require name lookup or access control (unlike by-name mechanisms), and it works even for bit fields (unlike mechanisms based on pointer-to-member values).

Note that this uses *expansion statements* as proposed by P1306 (and that is typical of practical uses of reflection). An expansion statement requires a compile-time range, which is why `data_members` must be a `constexpr` variable.

Schema generation

We can use this same pattern to generate SQL schemas from C++ classes. The implementation here mixes runtime SQL generation with static reflection, in order to demonstrate the interaction between these two features.

The entry point for the facility is a function template that takes a (standard layout) type parameter and writes the corresponding SQL CREATE TABLE statement.

```

template<StandardLayoutType T>
void create_table() {
    create_table_from_reflection<^T>();
}

```

This function simply delegates to a function parameterized by its reflection. Because reflection is expected to be an “advanced” feature, it might be desirable to hide it from user-facing interfaces. The SQL generating function template is shown below.

```

#include <meta>
using namespace meta = std::meta;

template<meta::info Class>
requires meta::is_class(Class)
void create_table_from_reflection() {

```

```

std::cout << "CREATE TABLE " << meta::name_of(Class) << "(\n";
constexpr auto members =
    meta::members_of(Class, is_non_static_data_member);
int size = members.size(), num = 0;
template for (constexpr meta::info member : members) {
    create_column<member>();
    if (++num != size)
        std::cout << ",\n";
}
std::cout << ");\n";
}

```

This function emits a CREATE TABLE statement for the name of the class, and “iterates” over the class’s data members — again using an *expansion statement* (P1306) — emitting column definitions for each (see below for `create_column`). We maintain the member count so that we can correctly insert commas into the output after each column.

Reflection facilities can only be used at compile time. Because this function mixes runtime code (`std::cout`) with static reflection (`meta::info`), we need to ensure that reflections do not “mix” with the runtime systems. We cannot, with this approach to generating SQL, pass the reflected class as a function argument, as that would leak the reflection — handle to an internal data structure that is only meaningful during translation — to run time. In other words, for mixed run-time/reflective algorithms reflection values must be passed as template arguments. We explore an alternative design of this algorithm in the following section.

Creating a column is straightforward: We serialize the member’s name and translate its C++ type into SQL.

```

template<meta::info Member>
requires meta::is_non_static_data_member(Member)
void create_column() {
    std::cout << meta::name_of(Member) << " ";
    std::cout << to_sql(meta::type_of(Member));
}

```

Finally, we need a facility to translate C++ types to SQL types. Here, we use a series of explicit specializations over reflections, with the generic case (i.e., primary template) triggering an instantiation error if used.

```

template<meta::info Type>
constexpr const char* to_sql() {
    static_assert(false, "no translation to SQL");
}

```

```
template<>
consteval const char* to_sql<^int>() {
    return "INTEGER";
}

template<>
consteval const char* to_sql<^float>() {
    return "FLOAT";
}
// etc.
```

Schema generation (take two)

The approach above mixes runtime SQL generation with static reflection: We call a function to print the schema to `std::cout`. An alternative approach is to synthesize the schema as a compile-time string, and then print the result later:

```
template<StandardLayoutType T>
consteval std::string create_table() {
    return create_table_from_reflection<^T>();
```

The string-generation code could be as follows:

```

#include <meta>
template<meta::info Class>
    requires meta::is_class(Class)
constexpr std::string create_table() {
    std::string result(""); // Assuming this should work
    result += "CREATE TABLE " + meta::name_of(Class) + "(\n";
    std::span<meta::info> members = data_members_of(^Class);
    int num = 0;
    for (meta::info member : members) {
        result += create_column(member);
        if (++num != members.size())
            result += ",\n";
    }
    result += ");\n";
    return result;
}

constexpr std::string void create_column(meta::info member) {
    std::string result("");
    result += meta::name_of(member) + " ";
    result += to_sql(meta::type_of(member));
    return result;
}

constexpr const char* to_sql(meta::info type) {
    static std::unordered_map<meta::info, const char*> types {
        {"^int", "INTEGER"},
        {"^float", "FLOAT"},
        // etc.
    };
    [[assert: meta::is_type(type)]];
    [[assert: types.count(type) != 0]]];
    return types.find(type)->second;
}

```

There are significant differences between this and the earlier example. In essence, this implementation looks like a normal program except that each function is a `constexpr` function dealing with reflection values. In other words, because the entire facility is expected to run at compile time, we don't have to use different constructs between the run-time and compile-time values in the implementation (e.g., expansion statements vs. loops); everything just looks like run-time code.

Template argument list assignment

In P0949R0, Peter Dimov proposes a facility to “assign” a list of template arguments for one template to another using:

```
mp_assign<ClassTmpl1<A1, A2, ...>, ClassTmpl2<B1, B2, ...>>
```

This is an alias for `ClassTmpl1<B1, B2, ...>`. The template arguments `An` and `Bn` are all type arguments. If the arguments of `mp_assign` are not of those forms, a substitution failure occurs.

Using the features proposed in this paper, we can implement this facility as follows:

```
#include <meta>
using std::meta::info;
using std::vector;
consteval info class_template_of(info inst) {
    using namespace std::meta;
    info tmpl = template_of(inst);
    if (is_class_template(inst) || is_invalid_reflection(tmpl)) {
        return tmpl;
    } else {
        return invalid_reflection("Not a class template instance");
    }
}

consteval vector<info> template_type_arguments_of(info inst) {
    using namespace std::meta;
    auto args = template_arguments_of(inst);
    for (auto arg: args) {
        if (is_invalid(arg) && args.size() == 1) {
            // template_arguments_of was invalid: Propagate the error.
            return vector<info>(args.begin(), args.end());
        } else if (!is_type(arg)) {
            // Not a type argument.
            return vector<info>{
                invalid_reflection("Not all arguments are types");
            }
        }
    }
    return vector<info>(args.begin(), args.end());
}
```

```

consteval info rf_assign(info inst1, info inst2) {
    using namespace std::meta;
    info tmpl1 = class_template_of(inst1);
    info tmpl2 = class_template_of(inst2);
    if (is_invalid(tmpl2)) return tmpl2;
    auto args1 = template_type_arguments_of(inst1);
    if (args1.size() == 1 && is_invalid(args1[0])) return args1;
    auto args2 = template_type_arguments_of(inst1);
    return substitute(tmpl1, args2);
}

```

If needed, `mp_assign` could be expressed in terms of `rf_assign`:

```

template<typename T1, typename T2>
using mp_assign = typename[:rf_assign(^T1, ^T2):];

```

Note that in our implementation of `rf_assign`, much of the code is dedicated to implementing the constraints of `mp_assign`. However, those constraints exist only because of two TMP limitations:

- 1) parameter packs cannot model mixed-kind template argument lists, and
- 2) template template parameters cannot accept function/variable templates.

In the reflection world we can easily lift those constraints, which produces the following simplified-yet-more-powerful implementation of `rf_assign`:

```

consteval info rf_assign(info inst1, info inst2) {
    using namespace std::meta;
    return substitute(template_of(inst1),
                      template_arguments_of(inst2));
}

```

Dealing more efficiently with parameter packs

Currently, parameter packs are generally dealt with through recursive template instantiation (i.e., a form of TMP, with all its disadvantages). With the set of features presented here, many interesting applications of packs can be expressed more directly and using fewer compilation resources. Here is a simple example:

```
#include <meta>
// Function taking an arbitrary number of arguments and returning a vector containing copies of the
// arguments that have the given type T.
template<typename T, typename ... Ts>
std::vector<T> select_values_of_type(Ts ... p) {
    std::vector<T> result{};
    template for (constexpr auto param:
                parameters_of(current_function())) {
        if (^T == type_of(param)) {
            result.push_back([:params[i]:]);
        }
    }
    return result;
}
```

Applying functions to all members

P0949r0 presents a TMP metafunction `get_all_data_members` aimed at collecting reflection information for all the data members of a class (not just the direct ones) using the facilities of the first reflection TS (N4818).

Unfortunately, `get_all_data_members` as presented in P0949r0 has a number of problems:

- It doesn't correctly use the TMP-based reflection API to access base classes (it looks like it treats a base class as a base class *type*). Fixing that is nontrivial.
- Its logic ignores virtual bases.
- The TMP-based reflection API doesn't deal well with bit fields (it relies on pointer-to-member constants, which cannot point to bit fields).

To address those shortcomings, we present a different interface with similar capabilities:

```
template<typename T, typename F>
void apply_to_all_data_members(T &&r_obj, F &&f);
// Invoke f(r_obj.x) for every accessible data member of r_obj, including
// those in base classes (and possibly hidden by more-derived member declarations).
```

With the facilities we have proposed in this paper, this can be implemented as follows.

```
#include <meta>
using std::meta::info;
```

```
// Convenience function to retrieve accessible nonstatic data members of a given class:
constexpr auto get_members(info classinfo) {
    return members_of(classinfo, is_nonstatic_data_member,
                      is_accessible);
}

// Convenience function to select nonvirtual bases and members.
constexpr auto is_not_virtual(info base_or_mem) {
    return !is_virtual(base_or_mem);
}

// Utility to get the reflection information for the types of base classes (rather than the base
// classes themselves) of a given class.
constexpr
auto get_base_types(info classtype, bool virtual_bases) {
    auto result = bases_of(classtype,
                           is_accessible,
                           virtual_bases ? is_virtual
                           : is_not_virtual);

    // Replace each base reflection by the reflection of its type.
    for (auto &info : result) {
        info = type_of(info);
    }
    return result;
}

template<typename T, typename F>
void apply_to_data_members(T *p_obj, F &f) {
    template for (constexpr auto member : get_members(^T)) {
        f(p_obj->[ :member:]);
    }
}
```

```

template<typename T, typename F>
void apply_to_base_data_members(T *p_obj, F &f,
                                bool virtual_bases,
                                bool skip_direct_members) {
    // Recursively traverse (depth-first) either the nonvirtual or virtual base classes (depending
    // on the virtual_bases flag). We do this by collecting the base class types and casting
    // the pointer one level up.
    auto type = ^T;
    template for (constexpr auto basetype :
                get_base_types(type, virtual_bases)) {
        apply_to_base_data_members<T, F>(
            static_cast<typename[ :basetype:]*>(p_obj), f,
            virtual_bases, /*skip_direct_members*/false);
    }
    if (!skip_direct_members) {
        // Now that the base classes have been traversed, handle the data members at this level.
        template for (constexpr auto member : get_members(type)) {
            f(p_obj->[ :member:]);
        }
    }
}

template<typename T, typename F>
void apply_to_all_data_members(T const &&r_obj, F &&f) {
    T const *p_obj = std::addressof(r_obj);
    apply_to_base_data_members<T, F>(
        p_obj, f, /*virtual_bases=*/true,
        /*skip_direct_members=*/true);
    apply_to_base_data_members<T, F>(
        p_obj, f, /*virtual_bases=*/false,
        /*skip_direct_members=*/false);
}

```

This implementation reads like ordinary C++ code. Every invocation instantiates three function templates, independently of how complex type `T` is (though the amount of code in each instantiation does depend on `T` because of the expansion statements).

This implementation still has a weakness, however: The notion of “accessibility” of bases and members is determined from the context of the implementation, not that of the call to `apply_to_all_data_members`. (The same limitation is imposed by the first Reflection TS.) We do not at this time propose to resolve that issue but we know of at least two ways to address it:

- more powerful code injection primitives, or
- introduce a reflection for “context”.

The second option would be less efficient since that context would have to be passed along as a template argument, which would cause each invocation of `apply_to_all_data_members` to have a distinct instantiation.

Alternatively, here is an implementation that doesn't use *expansion statements*. Instead, it relies on *fold-expressions*.

```
// Convenience function to retrieve accessible nonstatic data members of a given class:
consteval auto get_members(info classinfo) {
    return members_of(classinfo, is_nonstatic_data_member,
                      is_accessible);
}

// Convenience function to select nonvirtual bases and members.
consteval auto is_not_virtual(info base_or_mem) {
    return !is_virtual(base_or_mem);
}

// Utility to get the reflection information for the types of base classes (rather than the base
// classes themselves) of a given class.
consteval
auto get_base_types(info classtype, bool virtual_bases) {
    auto result = bases_of(classtype,
                           is_accessible,
                           virtual_bases ? is_virtual
                           : is_not_virtual);

    // Replace each base reflection by the reflection of its type.
    for (auto &info : result) {
        info = type_of(info);
    }
    return result;
};

template<typename T, typename F, std::meta::info ... members>
void apply_to_data_members(T *p_obj, F &f) {
    (void)(f(p_obj->[:members:])), ...); // Fold-expression.
}
```

```

template<typename T, typename F, std::meta::info ... classtypes>
void apply_to_base_data_members(T *p_obj, F &f,
                                bool virtual_bases,
                                bool skip_direct_members) {
    using namespace std::meta;
    // Use a fold-expression to recurse through given bases if needed.
    (apply_to_base_data_members<
        T, F, ...[:get_base_types(classtypes, virtual_bases):]...
    >(static_cast<typename(typeof(bases))*>(p_obj), f,
       virtual_bases,
       /*skip_direct_members=*/false), ...);
    if (!skip_direct_members) {
        // Use another fold-expression to handle the data members of each specified class type.
        (apply_to_data_members<
            T, F, ...[:get_members(classtypes):]...
        >(p_obj, f), ...);
    }
}

template<typename T, typename F>
void apply_to_all_data_members(T const &&r_obj, F &&f) {
    T const *p_obj = std::addressof(r_obj);
    apply_to_base_data_members<T, F, ^T>(
        p_obj, f, /*virtual_bases=*/true,
        /*skip_direct_members=*/true);
    apply_to_base_data_members<T, F, ^T>(
        p_obj, f, /*virtual_bases=*/false,
        /*skip_direct_members=*/false);
}

```

Clearly this is far less readable than the first version. It also involves more instantiations than the first version, but it is nonetheless more efficient than a pure TMP-based solution.

Appendix: Meta-library synopsis

This appendix briefly lists declarations for all the intrinsic meta-functions being worked on in the Lock3 Software implementation. As mentioned, these declarations are eventually meant to be brought into a program by including the standard header `<meta>`. In its current form the list differs slightly from the discussions in this paper because of implementation realities. We expect to harmonize the two over time.

Parameters to queries are named to represent the subset of `meta::info` values accepted by each function:

- `reflection` – accepts any value.

- `invalid` – accepts only an invalid reflection.
- `function` – accepts any value that designates a function.
- `variable` – accepts any value that designates a variable.
- `bitfield` – accepts any value that designates a bitfield.
- `type` – accepts any value that designates a type.
- `xxx_type` – accepts any value that designates a type of kind `xxx` (e.g., `enum_type`)
- `templ` – accepts any value that designates a template.
- `special` – accepts any value that designates a template specialization.
- `entity` – accepts any value that designates an entity.
- `parameter` – accepts any value that designates a function or template parameter.
- `expression` – accepts any value that designates an expression.
- `argument` – accepts any value that designates a function or template argument.
- `base` – accepts any value that designates a base class specifier.
- `mem` – accepts any value that designates a class member.
- `mem_function` – accepts any value that designates a member function.
- `spec_mem_function` – accepts any value that designates a special member function.
- `base_or_mem` – accepts any value that designates a base class specifier or a class member.

Some operations are polymorphic and accept combinations. For example, the parameter of `is_public` is `base_or_mem`. Operations accepting sequences are pluralized (e.g., `reflections`), meaning that the restriction applies to all elements of the sequence. Also, note that a function accepting a `declarator` will accept a function, variable, or bitfield.

Although all the declarations are being considered and worked on, some are not implemented or not tested. The list below highlights those declarations as follows:

Highlights for interfaces still requiring work	
Red	Not implemented, because of not-yet-resolved limitations
Orange	Not implemented, because of fixable limitations
Yellow	Not <i>yet</i> implemented
Green	Implemented, but not tested

The synopsis of the `<experimental/meta>` header is:

```
namespace std::meta {
    // Reflection type
    using info = decltype(^void);

    // Classification
    consteval bool is_invalid(info reflection);

    // Scope
    consteval bool is_local(info reflection);
    consteval bool is_class_member(info reflection);

    // Variables
    consteval bool is_variable(info reflection);
    consteval bool has_static_storage_duration(info variable);
    consteval bool has_thread_local_storage_duration(
        info variable);
    consteval bool has_automatic_storage_duration(
        info variable);

    // Functions
    consteval bool is_function(info reflection);
    consteval bool is_nothrow(info function);
    consteval bool has_ellipsis(info function);

    // Classes
    template<typename ...Args>
    consteval std::span<info> members_of(
        info class_type, Args ...filters);
    template<typename ...Args>
    consteval std::span<info> bases_of(
        info class_type, Args ...filters);

    // Classes
    consteval bool is_class(info reflection);
    consteval bool is_union(info reflection);
    consteval bool has_virtual_destructor(info class_type);
    consteval bool is_declared_class(info class_type);
    consteval bool is_declared_struct(info class_type);

    // Data Members
```

```
consteval bool is_data_member(info reflection);
consteval bool is_static_data_member(info reflection);
consteval bool is_nonstatic_data_member(info reflection);
consteval bool is_bit_field(info reflection);
consteval bool is mutable(info reflection);

// Member Functions
consteval bool is_member_function(info reflection);
consteval bool is_static_member_function(info reflection);
consteval bool is_nonstatic_member_function(info reflection);
consteval bool is_normal(info mem_function);
consteval bool is_conversion(info mem_function);
consteval bool is_override(info mem_function);
consteval bool is_override_specified(info mem_function);
consteval bool is_deleted(info mem_function);
consteval bool is_virtual(info mem_function);
consteval bool is_pure_virtual(info mem_function);

// Special Member Functions
consteval bool is_constructor(info reflection);
consteval bool is_default_constructor(info reflection);
consteval bool is_copy_constructor(info reflection);
consteval bool is_move_constructor(info reflection);
consteval bool is_copy_assignment_operator(info reflection);
consteval bool is_move_assignment_operator(info reflection);
consteval bool is_copy(info mem_function);
consteval bool is_move(info mem_function);
consteval bool is_destructor(info reflection);
consteval bool is_defaulted(info spec_mem_function);
consteval bool is_explicit(info spec_mem_function);

// Access
consteval bool has_access(info reflection);
consteval bool is_public(info base_or_mem);
consteval bool is_protected(info base_or_mem);
consteval bool is_private(info base_or_mem);
consteval bool has_default_access(info base_or_mem);

// Linkage
consteval bool has_linkage(info reflection);
consteval bool is_externally_linked(info reflection);
consteval bool is_internally_linked(info reflection);
```

```
// General purpose
consteval bool is_extern_specified(info reflection);
consteval bool is_inline(info reflection);
consteval bool is_inline_specified(info reflection);
consteval bool is_constexpr(info reflection);
consteval bool is_consteval(info reflection);
consteval bool is_final(info reflection);
consteval bool is_defined(info reflection);
consteval bool is_complete(info reflection);

// Namespaces
consteval bool is_namespace(info reflection);

// Aliases
consteval bool is_alias(info reflection);
consteval bool is_namespace_alias(info reflection);
consteval bool is_type_alias(info reflection);
consteval bool is_alias_template(info reflection);

// Enums
consteval bool is_enum(info reflection);
consteval bool is_unscoped_enum(info reflection);
consteval bool is_scoped_enum(info reflection);
consteval std::span<info> enumerators_of(info enum_type);

// Enumerator
consteval bool is_enumerator(info reflection);

// Templates
consteval bool is_template(info reflection);
consteval bool is_class_template(info reflection);
consteval bool is_function_template(info reflection);
consteval bool is_variable_template(info reflection);
consteval bool is_member_function_template(info reflection);
consteval bool is_static_member_function_template(reflection);
consteval bool is_nonstatic_member_function_template(
                           info reflection);
consteval bool is_constructor_template(info reflection);
consteval bool is_destructor_template(info reflection);
consteval bool is_concept(info reflection);

// Specializations
```

```
consteval bool is_specialization(info reflection);
consteval bool is_partial_specialization(info reflection);
consteval bool is_explicit_specialization(info reflection);
consteval bool is_implicit_instantiation(info reflection);
consteval bool is_explicit_instantiation(info reflection);

consteval info template_of(info special);
consteval bool has_template_arguments(info reflection);
consteval std::span<info> template_arguments_of(info special);
consteval info substitute(info templ, std::span<info> args);

// Base classes
consteval bool is_base_class(info reflection);
consteval bool is_direct_base_class(info reflection);
consteval bool is_virtual_base_class(info reflection);

// Parameters
consteval bool is_function_parameter(info reflection);
consteval bool is_template_parameter(info reflection);
consteval bool is_type_template_parameter(info reflection);
consteval bool is_nontype_template_parameter(info reflection);
consteval bool is_template_template_parameter(info reflection);
consteval bool has_default_argument(info parameter);

consteval std::span<info> parameters_of(info function_or_temp1);

// Types
consteval bool is_type(info reflection);
consteval bool is_fundamental_type(info type);
consteval bool has_fundamental_type(info reflection);
consteval bool is_arithmetic_type(info type);
consteval bool has_arithmetic_type(info reflection);
consteval bool is_scalar_type(info type);
consteval bool has_scalar_type(info reflection);
consteval bool is_object_type(info type);
consteval bool has_object_type(info reflection);
consteval bool is_compound_type(info type);
consteval bool has_compound_type(info reflection);
consteval bool is_function_type(info type);
consteval bool has_function_type(info reflection);
consteval bool is_class_type(info type);
consteval bool has_class_type(info reflection);
```

```
consteval bool is_union_type(info type);
consteval bool has_union_type(info reflection);
consteval bool is_enum_type(info type);
consteval bool has_enum_type(info type);
consteval bool is_unscoped_enum_type(info type);
consteval bool has_unscoped_enum_type(info reflection);
consteval bool is_scoped_enum_type(info type);
consteval bool has_scoped_enum_type(info reflection);
consteval bool is_void_type(info type);
consteval bool has_void_type(info reflection);
consteval bool is_null_pointer_type(info type);
consteval bool has_null_pointer_type(info reflection);
consteval bool is_integral_type(info type);
consteval bool has_integral_type(info reflection);
consteval bool is_floating_point_type(info type);
consteval bool has_floating_point_type(info reflection);
consteval bool is_array_type(info type);
consteval bool has_array_type(info reflection);
consteval bool is_pointer_type(info type);
consteval bool has_pointer_type(info reflection);
consteval bool is_reference_type(info type);
consteval bool has_reference_type(info reflection);
consteval bool is_lvalue_reference_type(info type);
consteval bool has_lvalue_reference_type(info reflection);
consteval bool is_rvalue_reference_type(info type);
consteval bool has_rvalue_reference_type(info reflection);
consteval bool is_member_pointer_type(info type);
consteval bool has_member_pointer_type(info reflection);
consteval bool is_member_object_pointer_type(info type);
consteval bool has_member_object_pointer_type(info reflection);
consteval bool is_member_function_pointer_type(info type);
consteval bool has_member_function_pointer_type(
                           info reflection);

consteval bool is_closure_type(info type);
consteval bool has_closure_type(info reflection);

// Type properties
consteval bool is_incomplete_type(info type);
consteval bool has_incomplete_type(info reflection);
consteval bool is_const_type(info type);
consteval bool has_const_type(info reflection);
consteval bool is_volatile_type(info type);
consteval bool has_volatile_type(info reflection);
```

```

consteval bool is_trivial_type(info type);
consteval bool has_trivial_type(info reflection);
consteval bool is_trivially_copyable_type(info type);
consteval bool has_trivially_copyable_type(info reflection);
consteval bool is_standard_layout_type(info type);
consteval bool has_standard_layout_type(info reflection);
consteval bool is_pod_type(info type);
consteval bool has_pod_type(info reflection);
consteval bool is_literal_type(info type);
consteval bool has_literal_type(info reflection);
consteval bool is_empty_type(info type);
consteval bool has_empty_type(info reflection);
consteval bool is_polymorphic_type(info type);
consteval bool has_polymorphic_type(info reflection);
consteval bool is_abstract_type(info type);
consteval bool has_abstract_type(info reflection);
consteval bool is_final_type(info type);
consteval bool has_final_type(info reflection);
consteval bool is_aggregate_type(info type);
consteval bool has_aggregate_type(info reflection);
consteval bool is_signed_type(info type);
consteval bool has_signed_type(info reflection);
consteval bool is_unsigned_type(info type);
consteval bool has_unsigned_type(info reflection);
consteval bool has_unique_object_representations(info type);
consteval bool has_type_with_unique_object_representations(
                           info reflection);

consteval std::size_t size_of(info reflection);
consteval std::size_t byte_size_of(info reflection);
consteval std::size_t bit_size_of(info reflection);
consteval std::size_t byte_offset_of(info reflection);
consteval std::size_t bit_offset_of(info reflection);
consteval std::size_t alignment_of(info reflection);
consteval std::size_t rank(info reflection);
consteval std::size_t extent(info reflection);

// Type operations
consteval bool is_constructible(
                           info reflection, std::span<info> arguments);
consteval bool is_trivially_constructible(
                           info reflection, std::span<info> arguments);
consteval bool is_nothrow_constructible(
                           info reflection, std::span<info> arguments);

```

```
        info reflection, std::span<info> arguments);  
consteval bool is_default_constructible_type(info type);  
consteval bool has_default_constructible_type(info reflection);  
consteval bool is_trivially_default_constructible_type(  
            info type);  
consteval bool has_trivially_default_constructible_type(  
            info reflection);  
consteval bool is_nothrow_default_constructible_type(info type);  
consteval bool has_nothrow_default_constructible_type(  
            info reflection);  
consteval bool is_copy_constructible_type(info type);  
consteval bool has_copy_constructible_type(info reflection);  
consteval bool is_trivially_copy_constructible_type(info type);  
consteval bool has_trivially_copy_constructible_type(  
            info reflection);  
consteval bool is_nothrow_copy_constructible_type(info type);  
consteval bool has_nothrow_copy_constructible_type(  
            info reflection);  
consteval bool is_move_constructible_type(info type);  
consteval bool has_move_constructible_type(info reflection);  
consteval bool is_trivially_move_constructible_type(info type);  
consteval bool has_trivially_move_constructible_type(  
            info reflection);  
consteval bool is_nothrow_move_constructible_type(info type);  
consteval bool has_nothrow_move_constructible_type(  
            info reflection);  
consteval bool is_assignable_type(info type,  
                                info assigned_type);  
consteval bool is_trivially_assignable_type(info type,  
                                         info assigned_type);  
consteval bool is_nothrow_assignable_type(info type,  
                                         info assigned_type);  
consteval bool is_copy_assignable_type(info type);  
consteval bool has_copy_assignable_type(info reflection);  
consteval bool is_trivially_copy_assignable_type(info type);  
consteval bool has_trivially_copy_assignable_type(  
            info reflection);  
consteval bool is_nothrow_copy_assignable_type(info type);  
consteval bool has_nothrow_copy_assignable_type(  
            info reflection);  
consteval bool is_move_assignable_type(info type);  
consteval bool has_move_assignable_type(info reflection);  
consteval bool is_trivially_move_assignable_type(info type);
```

```
consteval bool has_trivially_move_assignable_type(
    info reflection);
consteval bool is_nothrow_move_assignable_type(info type);
consteval bool has_nothrow_move_assignable_type(
    info reflection);
consteval bool is_destructible_type(info type);
consteval bool has_destructible_type(info reflection);
consteval bool is_trivially_destructible_type(info type);
consteval bool has_trivially_destructible_type(info reflection);
consteval bool is_nothrow_destructible_type(info type);
consteval bool has_nothrow_destructible_type(info reflection);

consteval bool is_swappable(info reflection);
consteval bool is_nothrow_swappable(info reflection);
consteval bool is_swappable_with(
    info reflection1, info reflection2);
consteval bool is_nothrow_swappable_with(
    info reflection1, info reflection2);

// Captures
consteval std::span<info> captures_of(info reflection);

consteval bool has_default_ref_capture(info reflection);
consteval bool has_default_copy_capture(info reflection);

consteval bool is_capture(info reflection);
consteval bool is_simple_capture(info reflection);
consteval bool is_ref_capture(info reflection);
consteval bool is_copy_capture(info reflection);
consteval bool is_explicit_capture(info reflection);
consteval bool is_init_capture(info reflection);
consteval bool has_captures(info reflection);

// Type relations
consteval bool is_same(info reflection1, info reflection2);
consteval bool is_base_of(info base_type, info derived_type);
consteval bool is_convertible(info from_type, info to_type);
consteval bool is_nothrow_convertible(
    info from_type, info to_type);

// Invocation
consteval bool is_invocable(
```

```
                info function, std::span<info> arguments);
consteval bool is_nothrow_invocable(
                info function, std::span<info> arguments);
consteval bool is_invocable_r(
                info function, std::span<info> arguments, info result);
consteval bool is_nothrow_invocable_r(
                info function, std::span<info> arguments, info result);

// Type transformation
consteval info remove_const(info type);
consteval info remove_volatile(info type);
consteval info remove_cv(info type);
consteval info add_const(info type);
consteval info add_volatile(info type);
consteval info add_cv(info type);
consteval info remove_reference(info type);
consteval info add_lvalue_reference(info type);
consteval info add_rvalue_reference(info type);
consteval info remove_pointer(info type);
consteval info add_pointer(info type);
consteval info remove_cvref(info type);
consteval info decay(info type);
consteval info make_signed(info type);
consteval info make_unsigned(info type);

consteval info aligned_storage(
                std::size_t length,
                std::size_t align = /* default-alignment */);
consteval info aligned_union(
                std::size_t length, std::span<info> types);

consteval info enable_if(bool cond, info type = ^void);

// Associated types
consteval info this_ref_type_of(info mem_function);
consteval info common_type(std::span<info> types);
consteval info underlying_type_of(info reflection);
consteval info invoke_result(
                info function, std::span<info> arguments);
consteval info type_of(info reflection);
consteval info return_type_of(info function);

// Associated reflections
```

```
consteval bool is_entity(info reflection);
consteval info entity_of(info reflection);
consteval info parent_of(info reflection);
consteval info definition_of(info reflection);

// Names
consteval bool is_named(info reflection);
consteval std::string_view name_of(info named);
consteval std::string_view display_name_of(info named);

// Expressions
consteval bool is_lvalue(info reflection);
consteval bool is_xvalue(info reflection);
consteval bool is_prvalue(info reflection);
consteval bool is_glvalue(info reflection);
consteval bool is_rvalue(info reflection);
}
```